

63 negations

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1024
1024 to 100
100 to 63
63

1024

DLL, SM

(Bounded) DLL:

Axioms:

$$A \vdash A$$

$$A \vdash \top$$

$$\perp \vdash A$$

$$A_0 \wedge A_1 \vdash A_i$$

$$A_i \vdash A_0 \vee A_1$$

$$A \wedge (B \vee C) \vdash (A \wedge B) \vee (A \wedge C)$$

Rules:

$$\frac{A \vdash B \quad A \vdash C}{A \vdash B \wedge C}$$

$$\frac{A \vdash C \quad B \vdash C}{A \vee B \vdash C}$$

$$\frac{A \vdash B \quad B \vdash C}{A \vdash C}$$

SM

$$\frac{A \vdash B}{-B \vdash -A}$$

Dualizing a sequent: swap premise/conclusion, \wedge/\vee , and \top/\perp .

Every axiom has a dual theorem. Every rule has a dual rule.

So every proof has a dual proof: a proof of the dual theorem.

Derived rule:

If $A \vdash B$ and $C()$ is a positive context, then $C(A) \vdash C(B)$.
If $C()$ is negative, then $C(B) \vdash C(A)$.

(Proof: induction on $C()$.)

1024

10 principles

Normality principles:

N1: $\top \vdash \perp$

N2: $\neg\top \vdash \perp$

N: N1 + N2

Antidistribution principles:

$$A1: \neg A \wedge \neg B \vdash \neg(A \vee B)$$

$$A2: \neg(A \wedge B) \vdash \neg A \vee \neg B$$

$$A: A1 + A2$$

Double negation principles:

D1: $A \vdash \neg \neg A$

D2: $\neg \neg A \vdash A$

D: D1 + D2

Minimal ex'ion principles:

$$x1: A \wedge \neg A \vdash \neg \top$$

$$x2: \neg \perp \vdash A \vee \neg A$$

$$x: x1 + x2$$

Recall: $\neg \top \vdash \neg B$, and $\neg B \vdash \neg \perp$.

Full ex'ion principles:

$$x^+1: A \wedge \neg A \vdash \perp$$

$$x^+2: \top \vdash A \vee \neg A$$

$$x^+: x^+1 + x^+2$$

The 1 principles and 2 principles are dual.

$2^{10} = 1024$ specifications.

How many distinct logics?

1024 to 100

1024 to 256

Clearly: $x_1 + N_2$ entails x^+1 .

Clearly: x^+1 entails x_1 .

Less clearly: x^+1 entails N_2 .

$$\begin{array}{l}
 \wedge R: \frac{-T \vdash T \quad -T \vdash -T}{-T \vdash T \wedge -T} \quad x+1: \frac{}{T \wedge -T \vdash \perp} \\
 \text{Cut: } \frac{}{-T \vdash \perp}
 \end{array}$$

Dually, X^+2 entails \aleph_1 .

No need for X^+ principles.
Down to $2^8 = 256$.

1024 to 100

256 to 100

D_i entails N_i :

$$\frac{\frac{\perp \vdash -T}{--T \vdash -\perp} \quad T \vdash --T}{T \vdash -\perp}$$

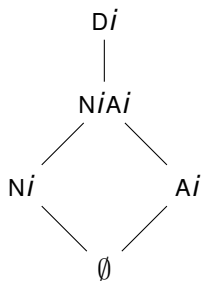
D_i entails A_i :

Part I

$$D2: \frac{\frac{-A \vdash -A \vee -B}{-(-A \vee -B) \vdash - - A}}{-(-A \vee -B) \vdash A}$$

Part II

$$D2: \frac{\frac{-(-A \vee -B) \vdash A \wedge B}{-(A \wedge B) \vdash - - (-A \vee -B)}}{-(A \wedge B) \vdash -A \vee -B}$$



This cuts $2^3 = 8$ down to 5.

256 was $(8 \times 2)^2$.

We reach $(5 \times 2)^2 = 100$.

100 to 63

Remaining entailments

That exhausts one-principle entailments.

Combinations remain.

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(Note: no use of distribution yet!)

Classical:

The pair of X^+ principles together entail all others.

So $X + N$ does too. Usual presentation of Boolean algebra.

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So $X + N$ does too. Usual presentation of Boolean algebra.

Of our 100, 15 are classical.

Full x brings full A.

Part I

$$\begin{array}{l}
 \text{Fiddling:} \quad \frac{\text{x2:} \quad \frac{}{\neg(A \wedge B) \vdash B \vee \neg B}}{\neg(A \wedge B) \vdash \neg(A \wedge B) \wedge (B \vee \neg B)}}{\text{Dist:} \quad \frac{\neg(A \wedge B) \vdash \neg(A \wedge B) \wedge (B \vee \neg B)}{\neg(A \wedge B) \vdash (\neg(A \wedge B) \wedge B) \vee (\neg(A \wedge B) \wedge \neg B)}} \\
 \text{DRule } (\wedge\text{elim}): \quad \frac{\neg(A \wedge B) \vdash (\neg(A \wedge B) \wedge B) \vee (\neg(A \wedge B) \wedge \neg B)}{\neg(A \wedge B) \vdash (\neg(A \wedge B) \wedge B) \vee \neg B}
 \end{array}$$

Part II

$$\begin{array}{l}
 \text{Fiddling:} \quad \frac{\text{x2:} \quad \frac{}{\neg(A \wedge B) \vdash A \vee \neg A}}{\neg(A \wedge B) \wedge B \vdash (\neg(A \wedge B) \wedge B) \wedge (A \vee \neg A)}}{\text{Dist:} \quad \frac{\neg(A \wedge B) \wedge B \vdash (\neg(A \wedge B) \wedge B) \wedge (A \vee \neg A)}{\neg(A \wedge B) \wedge B \vdash (\neg(A \wedge B) \wedge B \wedge A) \vee (\neg(A \wedge B) \wedge B \wedge \neg A)}} \\
 \text{DR } (\wedge\text{E}): \quad \frac{\neg(A \wedge B) \wedge B \vdash (\neg(A \wedge B) \wedge B \wedge A) \vee (\neg(A \wedge B) \wedge B \wedge \neg A)}{\neg(A \wedge B) \wedge B \vdash (\neg(A \wedge B) \wedge B \wedge A) \vee \neg A} \\
 \text{DR } (\text{x1, fiddling}): \quad \frac{\neg(A \wedge B) \wedge B \vdash (\neg(A \wedge B) \wedge B \wedge A) \vee \neg A}{\text{Fiddling:} \quad \frac{\neg(A \wedge B) \wedge B \vdash \neg A \vee \neg A}{\neg(A \wedge B) \wedge B \vdash \neg A}}
 \end{array}$$

Full x plus Ni brings Di .

$$\text{Dist: } \frac{}{\frac{}{\neg\neg A \wedge (A \vee \neg A) \vdash (\neg\neg A \wedge A) \vee (\neg\neg A \wedge \neg A)}}{\quad}$$

D Rule ($x^{+1} = x1N2$):

$$\perp\text{-drop: } \frac{\neg\neg A \wedge (A \vee \neg A) \vdash (\neg\neg A \wedge A) \vee \perp}{\neg\neg A \wedge (A \vee \neg A) \vdash \neg\neg A \wedge A}$$

$$\wedge\text{-elim: } \frac{\neg\neg A \wedge (A \vee \neg A) \vdash \neg\neg A \wedge A}{\neg\neg A \wedge (A \vee \neg A) \vdash A}$$

Full N plus x_i and $A\bar{i}$ brings D_i .

$$\begin{array}{l}
 \text{Dist:} \\
 \text{Fiddling (DRule, x1, N2):} \\
 \text{DRule (A2):} \\
 \text{DRule (x1,N2):} \\
 \text{DRule (N1):} \\
 \text{T-drop:}
 \end{array}
 \frac{
 \frac{
 \frac{
 A \wedge (-A \vee \neg \neg A) \vdash (A \wedge \neg A) \vee (A \wedge \neg \neg A)
 }{
 A \wedge (-A \vee \neg \neg A) \vdash \neg \neg A
 }{
 A \wedge \neg(A \wedge \neg A) \vdash \neg \neg A
 }{
 A \wedge \neg \perp \vdash \neg \neg A
 }{
 A \wedge \top \vdash \neg \neg A
 }{
 A \vdash \neg \neg A
 }
 }{
 }
 }{
 }
 }{
 }
 }{
 }
 }{
 }$$

Di plus $x\bar{i}$ brings xi .

$$x2: \frac{}{-\perp \vdash A \vee -A}$$

$$SM: \frac{}{-(A \vee -A) \vdash - - \perp}$$

$$DR(N1): \frac{}{-(A \vee -A) \vdash -\top}$$

$$\begin{array}{l} \text{Fiddling:} \\ \text{Cut:} \end{array} \frac{\frac{A \vdash - - A}{-A \wedge A \vdash -A \wedge - - A} \quad A1: \frac{}{-A \wedge - - A \vdash -(A \vee -A)}}{-A \wedge A \vdash -(A \vee -A)}$$

x_i plus \bar{n}_i plus A_i brings A_i .

Part I

$$\begin{array}{l}
 \text{Fiddling:} \quad \frac{x_1: \quad \frac{-A \wedge \neg \neg A \triangleright \neg(A \vee B)}{\quad}}{(-A \wedge \neg \neg A) \vee (-A \wedge \neg(A \vee B)) \triangleright \neg(A \vee B)} \\
 \text{Dist:} \quad \frac{\quad}{-A \wedge (\neg \neg A \vee \neg(A \vee B)) \triangleright \neg(A \vee B)} \\
 \text{DRule (A2):} \quad \frac{\quad}{-A \wedge \neg(\neg A \wedge (A \vee B)) \triangleright \neg(A \vee B)}
 \end{array}$$

Part II:

$$\begin{array}{l}
 \text{ECQ (x1 + N2):} \quad \frac{\quad}{A \wedge \neg A \triangleright B} \\
 \text{Fiddling:} \quad \frac{\quad}{(-A \wedge A) \vee (-A \wedge B) \triangleright B} \\
 \text{Cut (Dist):} \quad \frac{\quad}{-A \wedge (A \vee B) \triangleright B}
 \end{array}$$

That's it!

Down to 63.

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Some interesting critters

CL is NX (= DX).

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FDE is D.

Preminimal logic (the logic of compatibility frames) is $N1A1$.

Dual premin (the logic of exhaustiveness frames) is $N2A2$.

Ockham logic (the logic of the Routley star) is NA.

$D1X1$ is the strongest logic here sound for minimal logic,
 $N2D1X1$ for intuitionist.

$D2X2$ is the strongest logic here sound for dual minimal logic,
 $N1D2X2$ for dual intuitionist.

$D1X1$ is the strongest logic here sound for minimal logic,
 $N2D1X1$ for intuitionist.

$D2X2$ is the strongest logic here sound for dual minimal logic,
 $N1D2X2$ for dual intuitionist.

Are they these logics?

Other than CL, there are two “attractive” logics among the 100:

L1:

$$X2D1 = XN1 = XA2D1$$

L2:

$$X1D2 = XN2 = XA1D2$$

Anybody recognize these?

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Organized by x

No x: 25 logics

\cancel{x}	0	N1	A1	N1A1	D1
0	*	*	*	*(Pre)	*(Quas)
N2	*	*	*	*	*
A2	*	*	*	*	*
N2A2	*(DPre)	*	*	*(Ock)	*
D2	*(DQuas)	*	*	*	*(FDE)

x1: 17 logics

x1	0	N1	A1	N1A1	D1
0	*	*	*	*	* (Min)
N2	*	*	*	*	* (Int)
A2	*	*	*	*	*
N2A2	A1	D1	*	D1	*
D2	X2A1 (L2)	X2A1 (C)	X2 (L2)	X2 (C)	X2 (C)

x2: 17 logics

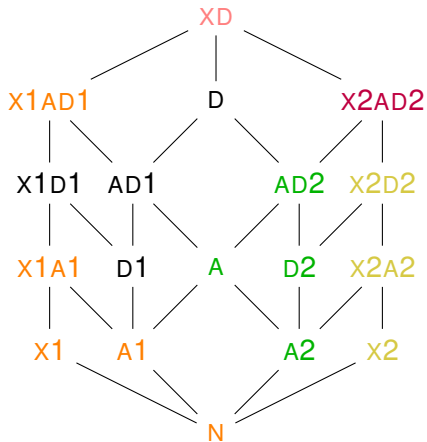
x2	0	N1	A1	N1A1	D1
0	*	*	*	A2	X1A2 (L1)
N2	*	*	*	D2	X1A2 (C)
A2	*	*	*	*	X1 (L1)
N2A2	*	*	*	D2	X1 (C)
D2	* (DMin)	* (DInt)	*	*	X1 (C)

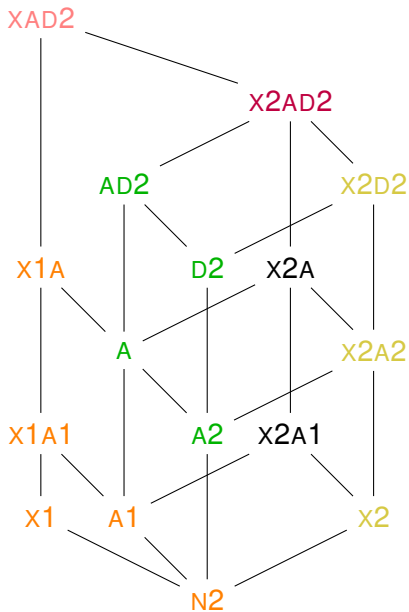
x: 4 logics

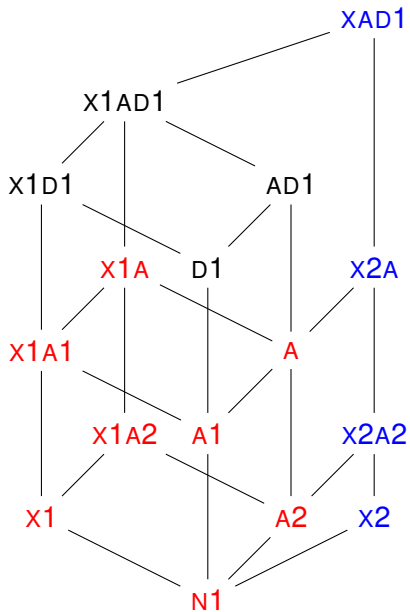
X	0	N1	A1	N1A1	D1
0	A	D1 (L1)	A2	D1 (L1)	A2 (L1)
N2	D2 (L2)	D (C)	D2 (L2)	D (C)	D2 (C)
A2	A1	D1 (L1)	*	D1 (L1)	* (L1)
N2A2	D2 (L2)	D (C)	D2 (L2)	D (C)	D2 (C)
D2	A1 (L2)	D1 (C)	* (L2)	D1 (C)	* (C)

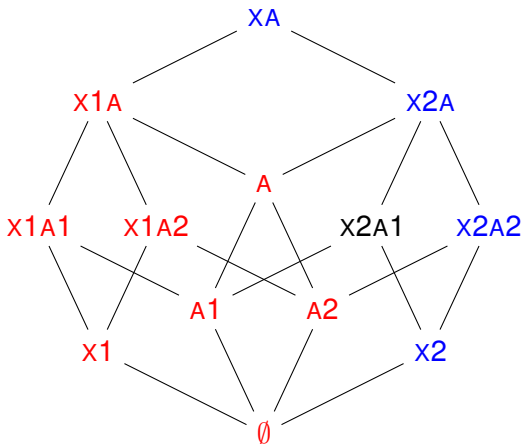
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Organized by N









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Next steps

Consider the rules of antilogism:

Antilogisms

$$\frac{A \wedge B \vdash C}{A \wedge \neg C \vdash \neg B}$$

$$\frac{A \vdash B \vee C}{\neg B \vdash \neg A \vee C}$$

$$\frac{A \wedge B \vdash C \vee D}{A \wedge \neg C \vdash \neg B \vee D}$$

Now how many?

Which of these critters can coexist?

Which of these critters can coexist?

How do they relate to other critters?